

On the Structure of Expected Payoff Sets in Multi-Objective Markov Decision Processes

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CFV – October 11, 2024

Random strategies and multiple objectives

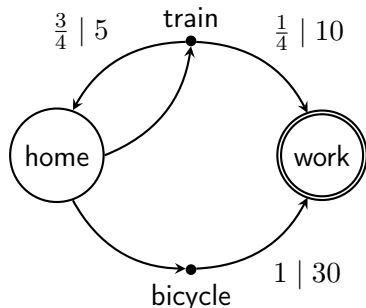
- We study **Markov decision processes** with **multiple payoffs**.
- In general, the satisfaction of multi-objective queries requires **randomised strategies**.

Main questions

- What is the relationship between expected payoffs of **pure strategies** and expected payoffs of **general strategies**?
- What **type of randomisation** do we need for multi-objective queries?

→ **Goal**: results for the **broadest possible class of payoffs**.

Markov decision processes



Markov decision process \mathcal{M}

- **Finite** state space S
- **Finite** action space A
- **Randomised** transitions

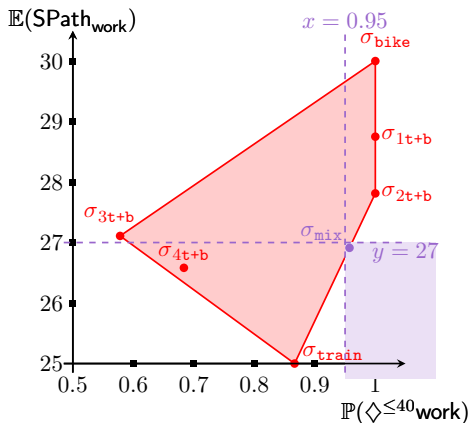
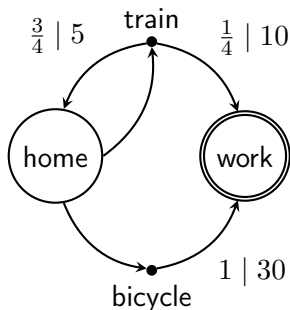
Plays are sequences in $(SA)^\omega$ coherent with transitions.

- A **strategy** is a function $\sigma: (SA)^*S \rightarrow \mathcal{D}(A)$
- A strategy σ is **pure** if it is not randomised.
- A strategy σ and initial state s induce a **distribution** \mathbb{P}_s^σ over **plays**.
- A **payoff** is a measurable function $f: \text{Plays}(\mathcal{M}) \rightarrow \bar{\mathbb{R}}$.
- We let $\mathbb{E}_s^\sigma(f) = \int_{\pi \in \text{Plays}(\mathcal{M})} f(\pi) d\mathbb{P}_s^\sigma(\pi)$.

Multi-objective Markov decision processes

We consider **two goals**:

- reaching work under 40 minutes with **high probability**;
- minimising the **expectancy** of the time to reach work.



What are good payoffs?

To provide formal results, we need to **constrain considered payoffs**.

$\rightsquigarrow \mathbb{E}_s^\sigma(f)$ should be well-defined **for all strategies**.

Good payoff functions

Three types of **good** payoffs:

- **non-negative** payoffs: $f \geq 0$;
- **non-positive** payoffs: $f \leq 0$;
- **universally integrable** payoffs: $\mathbb{E}_s^\sigma(|f|) \in \mathbb{R}$ for all strategies σ and all $s \in S$.

For a **multi-dimensional payoff** $\bar{f} = (f_1, \dots, f_d)$ and $s \in S$, we let:

- $\text{Pay}_s(\bar{f}) = \{\mathbb{E}_s^\sigma(\bar{f}) \mid \sigma \text{ strategy}\}$;
- $\text{Pay}_s^{\text{pure}}(\bar{f}) = \{\mathbb{E}_s^\sigma(\bar{f}) \mid \sigma \text{ pure strategy}\}$.

Universally integrable payoffs

In the introductory example, we had $\text{Pay}_{\text{home}}(\bar{f}) = \text{conv}(\text{Pay}_{\text{home}}^{\text{pure}}(\bar{f}))$.

When does this generalise?

Theorem

Let $\bar{f} = (f_1, \dots, f_d)$ be **universally integrable**. Then, for all states s ,

$$\text{Pay}_s(\bar{f}) = \text{conv}(\text{Pay}_s^{\text{pure}}(\bar{f})).$$

In particular, to match the expected payoff of any strategy, it suffices to:

- **mix $d + 1$ pure strategies;**
- *consider strategies use **randomisation at most d** along any play.*

Sequel: proof of a weaker result

If \bar{f} is universally integrable, then $\text{cl}(\text{Pay}_s(\bar{f})) = \text{cl}(\text{conv}(\text{Pay}_s^{\text{pure}}(\bar{f})))$.

Universally integrable payoffs

A simpler proof

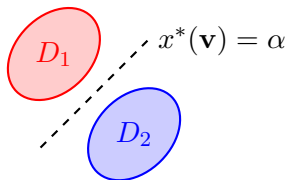
Non-direct inclusion: $\text{Pay}_s(\bar{f}) \subseteq \text{cl}(\text{conv}(\text{Pay}_s^{\text{pure}}(\bar{f})))$.

Let σ be a strategy and $\mathbf{q} = \mathbb{E}_s^\sigma(\bar{f})$. Assume $\mathbf{q} \notin \text{cl}(\text{conv}(\text{Pay}_s^{\text{pure}}(\bar{f})))$.

Main idea: reduction to a **one-dimensional** payoff.

Theorem (Hyperplane separation theorem)

Let $D_1, D_2 \subseteq \mathbb{R}^d$ be **disjoint convex** sets. If D_1 is **closed** and D_2 is **compact**, then there exists a **linear form** $x^*: \mathbb{R}^d \rightarrow \mathbb{R}$ and $\varepsilon > 0$ such that for all $\mathbf{p}_1 \in D_1$ and $\mathbf{p}_2 \in D_2$, $x^*(\mathbf{p}_1) + \varepsilon < x^*(\mathbf{p}_2)$.



Universally integrable payoffs

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Main idea: reduction to a **one-dimensional** payoff.

- There exists a linear form x^* such that, for all **pure strategies** τ ,

$$x^*(\mathbb{E}_s^\tau(\bar{f})) < x^*(\mathbf{q}).$$

- By linearity, we obtain that for all pure strategies τ ,

$$\mathbb{E}_s^\tau(x^*(\bar{f})) < \mathbb{E}_s^\sigma(x^*(\bar{f})).$$

Lemma

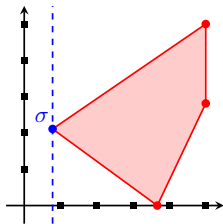
Let f be **universally integrable**. For all strategies σ , there exists a **pure strategy** τ such that $\mathbb{E}_s^\sigma(f) \leq \mathbb{E}_s^\tau(f)$.

Universally integrable payoffs

Compact case

What happens if $\text{Pay}_s(\bar{f})$ is compact ?

- The argument can be adapted if $\text{Pay}_s(\bar{f})$ is **polyhedral**.
- **However, good hyperplanes** do not generally exist for all extreme points despite $\text{Pay}_s(\bar{f}) = \text{conv}(\text{extr}(\text{Pay}_s(\bar{f})))$.



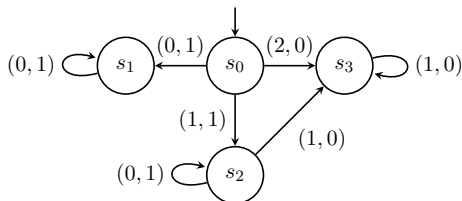
- Consider a **vertex** \mathbf{q} obtained by σ .
- There is a **hyperplane** intersecting $\text{Pay}_s(\bar{f})$ only at \mathbf{q} .
- There exists a **linear form** x^* such that σ is **optimal** from s for $x^* \circ \bar{f}$.
 $\rightsquigarrow \mathbf{q} \in \text{Pay}_s^{\text{pure}}(\bar{f})$

Universally integrable payoffs

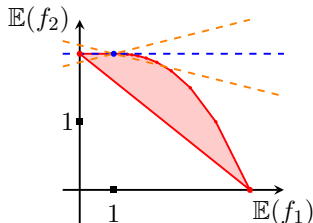
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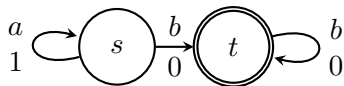


$$f_1(s_0 a_0 s_1 \dots) = \sum_{\ell=0}^{\infty} \left(\frac{3}{4}\right)^\ell w_1(a_\ell); \quad f_2(s_0 a_0 s_1 \dots) = \sum_{\ell=0}^{\infty} \left(\frac{1}{2}\right)^\ell w_2(a_\ell).$$



Beyond universally integrable payoffs

Example



Payoffs

- 1 reaching $t \rightsquigarrow f_1 = \mathbb{1}_{\diamond t}$;
- 2 sum of weights $\rightsquigarrow f_2 = \sum_{\ell=0}^{\infty} w(c_{\ell})$.

- $\mathbb{E}_s^{\sigma_a}(f_2) = +\infty \implies f_2$ is **not universally integrable**.
- $\text{Pay}_s^{\text{pure}}(\bar{f}) = \{(0, +\infty)\} \cup \{(1, \ell) \mid \ell \in \mathbb{N}\}$.
 $\implies \text{conv}(\text{Pay}_s^{\text{pure}}(\bar{f})) = (\{1\} \times \mathbb{R}_{\geq 0}) \cup ([0, 1[\times \{+\infty\})$.
- We have $(1, +\infty) \in \text{Pay}_s(\bar{f})$ via σ such that for all $\ell \in \mathbb{N}$:

$$\sigma(s(as)^{\ell})(a) = \begin{cases} \frac{1}{2} & \text{if } \ell \in 2^{\mathbb{N}} \\ 1 & \text{if } \ell \notin 2^{\mathbb{N}} \end{cases}$$

→ **The theorem does not generalise.**

Beyond universally integrable payoffs

Theorem

Let $\bar{f} = (f_1, \dots, f_d)$ be a good payoff and $s \in S$.

For all strategies σ , all $\varepsilon > 0$ and all $M \in \mathbb{R}$, there exist finitely many pure strategies τ_1, \dots, τ_n and coefficients $\alpha_1, \dots, \alpha_n \in [0, 1]$ such that

$\sum_m^n \alpha_m = 1$ and for all $1 \leq j \leq d$:

- if $\mathbb{E}_s^\sigma(f_j) = +\infty$, then $\sum_{m=1}^n \alpha_m \mathbb{E}_s^{\tau_m}(f_j) \geq M$,
- if $\mathbb{E}_s^\sigma(f_j) = -\infty$, then $\sum_{m=1}^n \alpha_m \mathbb{E}_s^{\tau_m}(f_j) \leq -M$, and,
- otherwise, if $\mathbb{E}_s^\sigma(f_j) \in \mathbb{R}$,
 $\mathbb{E}_s^\sigma(f_j) - \varepsilon \leq \sum_{m=1}^n \alpha_m \mathbb{E}_s^{\tau_m}(f_j) \leq \mathbb{E}_s^\sigma(f_j) + \varepsilon$.

Informally, we have

$$\text{cl}(\text{Pay}_s(\bar{f})) = \text{cl}(\text{conv}(\text{Pay}_s^{\text{pure}}(\bar{f}))).$$

Thank you for your attention !

Universally integrable payoffs

General argument: sketch

Proof goal

For all strategies σ , $\mathbf{q} = \mathbb{E}_s^\sigma(\bar{f}) \in \text{conv}(\text{Pay}_s^{\text{pure}}(\mathbf{q}))$.

- Construct linear map $L_{\mathbf{q}}$ such that:
 - a σ is **lexicographically optimal** for $L_{\mathbf{q}} \circ \bar{f}$;
 - b $\mathbf{q} \in \text{ri}(\text{Pay}_s(\bar{f}) \cap V)$ where $V = \{\mathbf{p} \in \mathbb{R}^d \mid L_{\mathbf{q}}(\mathbf{p}) = L_{\mathbf{q}}(\mathbf{q})\}$.
- Show that $\text{ri}(\text{Pay}_s(\bar{f}) \cap V) = \text{ri}(\text{conv}(\text{Pay}_s^{\text{pure}}(\bar{f})) \cap V)$, i.e.,
$$\text{cl}(\text{Pay}_s(\bar{f}) \cap V) = \text{cl}(\text{conv}(\text{Pay}_s^{\text{pure}}(\bar{f})) \cap V)$$

Key lemma

If \bar{f} is universally integrable, then for all strategies σ and all $s \in S$, there exists a pure strategy τ such that

$$\mathbb{E}_s^\sigma(\bar{f}) \leq_{\text{lex}} \mathbb{E}_s^\tau(\bar{f}).$$

A set of expected payoffs that is not closed

For $j \in \{1, 2\}$, we consider the payoff f_j such that, for all plays $s_0 a_0 s_1 \dots$,

$$f_j(s_0 a_0 s_1 \dots) = \mathbb{1}_{\text{Reach}(\{t\})(s_0 a_0 s_1 \dots)} \cdot \sum_{\ell=0}^{\infty} \left(\frac{3}{4}\right)^\ell w_j(a_\ell).$$

